

# Mapping Flows in Bipartite Networks

NetSci 2020

<https://arxiv.org/abs/2007.01666>

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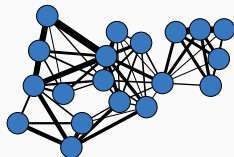
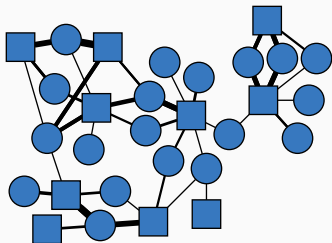
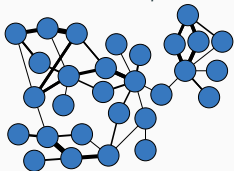
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## The Problem: Community Detection in Bipartite Networks

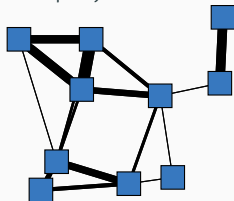
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# The Problem: Community Detection in Bipartite Networks

Treat as unipartite



Unipartite  
projections



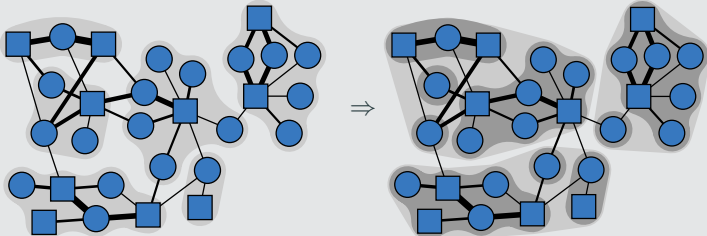
# The Problem: Community Detection in Bipartite Networks

## Our Solution

Teach the map equation to recognise bipartite networks.

## The Benefits

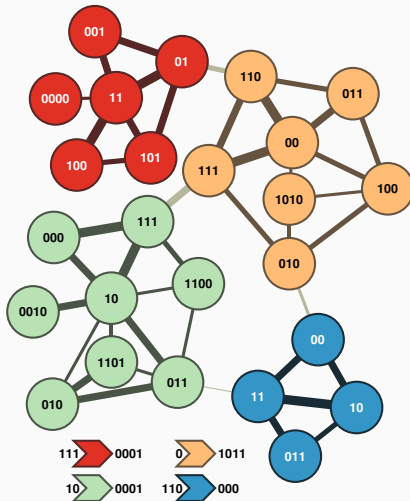
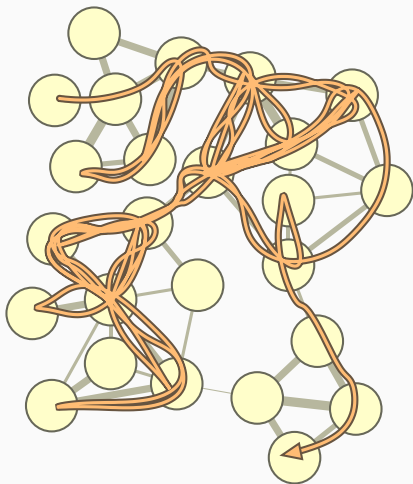
- We use all available data efficiently
- This improves the compression and we find more regularities
- We increase the resolution and explore different scales



## Quick Overview: The Map Equation Framework

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# The Map Equation Framework



$$L(M) = qH(Q) + \sum_{m \in M} p_m H(P_m)$$

Our Solution:  
The Bipartite Map Equation

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# The Bipartite Map Equation

## Idea

Reflect the bipartite network structure in the coding scheme.

## Key insight

Random walks must alternate between node types.

- network flow is divided evenly between node types,
- one node type is visited in even steps, the other in odd steps,
- there are two random processes:
  - X: the current node
  - Y: the current node type
  - then we can use Bayes' rule:  $H(X|Y) = H(X) - H(Y) + H(Y|X)$

$$L(M_1) = \underbrace{H(\mathcal{P})}_{H(X)} = \underbrace{1}_{H(Y)} + \underbrace{\frac{1}{2}H(\mathcal{P}^L) + \frac{1}{2}H(\mathcal{P}^R)}_{H(X|Y)}$$

- plug this into the map equation...



# The Bipartite Map Equation

$$L(M) = qH(Q) + \sum_{m \in M} p_m H(P_m)$$

↓

$$\begin{aligned} L_B(M) &= q^L H(Q^L) + \sum_{m \in M} p_m^L H(P_m^L) \\ &\quad + q^R H(Q^R) + \sum_{m \in M} p_m^R H(P_m^R) \end{aligned}$$

## Problem

In sparse networks, knowing the node type comes close to knowing the exact node  $\rightarrow$  encoding close to the entropy rate of the Markov process without identifying modular structure.

A key ingredient of the map equation is forgetting the right amount!

Our Solution 2.0:  
The Bipartite Map Equation  
with Varying Node-Type Memory

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# The Bipartite Map Equation with Varying Node-Type Memory

## Idea

Reflect the bipartite network structure in the coding scheme  
and forget node types at rate  $\alpha$ .

## Then

- Random walks still alternate between node types
- Random walker “confuses” node types with probability  $\alpha$   
→ new visit rates:  $p_u \rightsquigarrow p_u^\alpha = ((1 - \alpha) p_u, \alpha p_u)$  for left nodes  $u$   
 $p_v \rightsquigarrow p_v^\alpha = (\alpha p_v, (1 - \alpha) p_v)$  for right nodes  $v$
- Network flow is still evenly divided between node types!
- Before:  $H(Y|X) = 0$   
Now:  $H(Y|X) = H(\{\alpha, 1 - \alpha\}) = H_\alpha$

$$L(M_1) = \underbrace{H(\mathcal{P}_1)}_{H(X)} = \underbrace{1}_{H(Y)} - \underbrace{H_\alpha}_{H(Y|X)} + \underbrace{H(\mathcal{P}_1^\alpha)}_{H(X|Y)}$$

→ plug this into the map equation...

# The Bipartite Map Equation with Varying Node-Type Memory

$$L(M) = qH(Q) + \sum_{m \in M} p_m H(P_m)$$



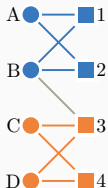
## The Bipartite Map Equation with Varying Node-Type Memory

$$L^\alpha(M) = q^\alpha H(Q^\alpha) + \sum_{m \in M} p_m^\alpha H(P_m^\alpha)$$

For  $\alpha = \frac{1}{2}$ , we recover the standard map equation!

# The Bipartite Map Equation with Varying Node-Type Memory

(a)



(b)

$$L(M) = \underbrace{\begin{bmatrix} \text{blue} & \text{orange} \end{bmatrix}}_{q = \frac{1}{9}} H(Q) = 1 \quad + \quad \begin{matrix} p_1 = \frac{5}{9} \\ \begin{bmatrix} \text{A} & \text{B} & 1 & 2 & e \end{bmatrix} H(P_1) = 2.25 \\ \begin{bmatrix} \text{C} & \text{D} & 3 & 4 & e \end{bmatrix} H(P_2) = 2.25 \\ p_2 = \frac{5}{9} \end{matrix} = 2.61$$

Index Level
Module Level

(c)

$$L^{0.1}(M) = \underbrace{\begin{bmatrix} \text{blue} & \text{orange} \end{bmatrix}}_{q = (\frac{1}{18}, \frac{1}{18})} H(Q^{0.1}) = \begin{Bmatrix} 0.47 \\ 0.47 \end{Bmatrix} + \begin{matrix} p_1 = (\frac{5}{18}, \frac{5}{18}) \\ \begin{bmatrix} \text{A} & \text{B} & \text{III} \\ \text{I} & 2 & e \end{bmatrix} H(P_1^{0.1}) = \begin{Bmatrix} 1.50 \\ 1.94 \end{Bmatrix} \\ \begin{bmatrix} \text{C} & \text{D} & e \\ \text{III} & 3 & 4 \end{bmatrix} H(P_2^{0.1}) = \begin{Bmatrix} 1.94 \\ 1.50 \end{Bmatrix} \\ p_2 = (\frac{5}{18}, \frac{5}{18}) \end{matrix} = 1.96$$

## Results

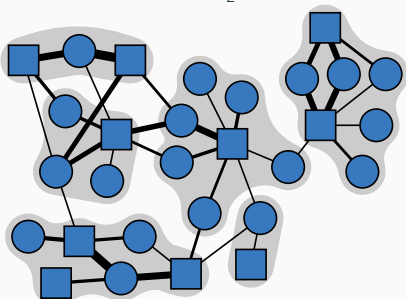
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# Results

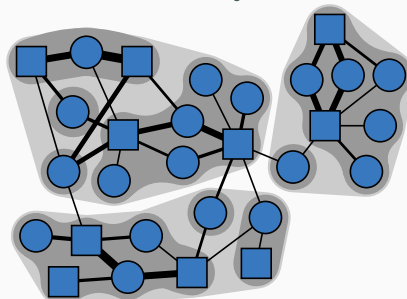
## The Bipartite Map Equation with Varying Node-Type Memory

$$L^{\alpha}(M) = q^{\alpha} H(\mathcal{Q}^{\alpha}) + \sum_{m \in M} p_m^{\alpha} H(\mathcal{P}_m^{\alpha})$$

$$\alpha = \frac{1}{2}$$



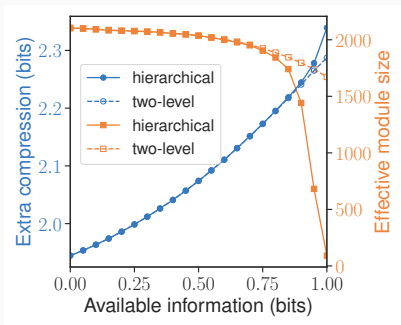
$$\alpha = \frac{1}{6}$$



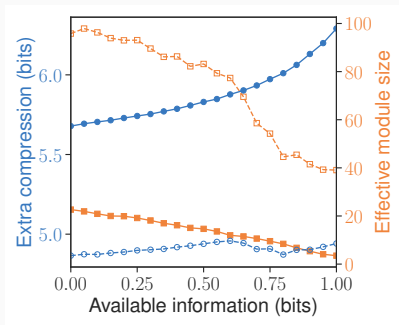
$\Rightarrow$  we can change the resolution and explore different scales!

# Results

(a) Last.fm user-song network.



(b) IMDb actor-movie network.



Extra compression: compared to one-module partition

Available information (bits) =  $1 - H_\alpha$

Effective module size =  $2^{H(S)}$



## Conclusion

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## Problem

Most community detection methods address unipartite networks.

## Idea

Extend the map equation framework to reflect the regularities in bipartite networks.

## Benefits

We use node-type information efficiently and increase the compression. Further, by adjusting the resolution, we explore community structures at different scales.

<https://arxiv.org/abs/2007.01666>

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